# Differentially Private CountSketch

#### Improved utility analysis

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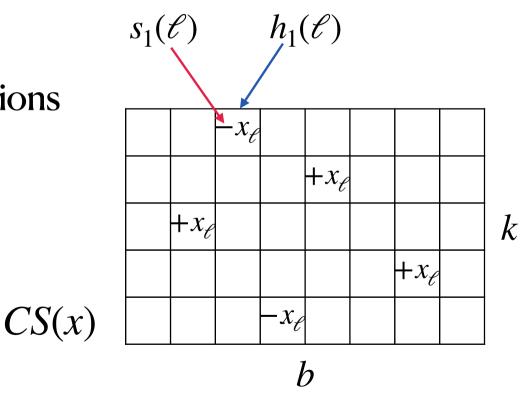


### **CountSketch**

[Charikar, Chen, Farach-Colton 2002]

- Linear sketch,  $CS : \mathbf{R}^d \to \mathbf{R}^{k \times b}$
- Defined using random hash functions  $h_1, \dots, h_k : [d] \rightarrow [b]$  $s_1, \dots, s_k : [d] \rightarrow \{-1, +1\}$

**This talk**: Assume hash functions are *fully* independent

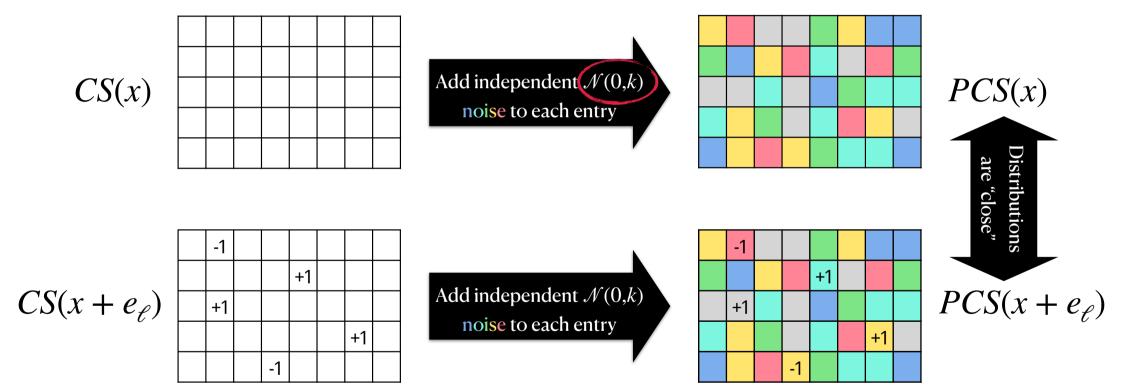


#### **CountSketch estimator**

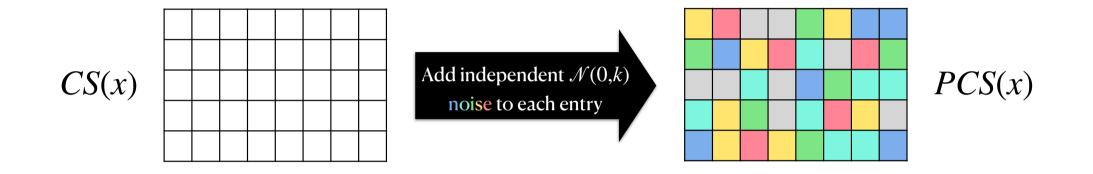
- Simple estimators:  $s_1(\ell)CS(x)_{1,h_1(\ell)}, \ldots, s_k(\ell)CS(x)_{k,h_k(\ell)}$
- Median estimator:  $\hat{x}_{\ell} = \text{median}(s_i(\ell)CS(x)_{i,h_i(\ell)} \mid i \in [k])$

**Theorem** (Minton & Price, 2014) For every  $\alpha \in [0, 1]$  and  $\Delta = ||\operatorname{tail}_b(x)||_2/\sqrt{b}$ ,  $\Pr[|\hat{x}_{\ell} - x_{\ell}| > \alpha \Delta] < 2 \exp(-\Omega(\alpha^2 k))$ ,  $\Delta$  is "maximum error of CountSketch"

## Making CountSketch differentially private



#### **Estimation from Private CountSketch**



 $\hat{x}_{\ell} = \text{median}(s_i(\ell)CS(x)_{i,h_i(\ell)} \mid i \in [k]) \qquad \bar{x}_{\ell} = \text{median}(s_i(\ell)PCS(x)_{i,h_i(\ell)} \mid i \in [k])$ 

**The question**: How much worse is the private estimator  $\bar{x}_{\ell}$  compared to  $\hat{x}_{l}$ ?

## Our result

**Theorem** For every  $\alpha \in [0, 1]$  and  $\Delta = ||\operatorname{tail}_b(x)||_2 / \sqrt{b}$ ,  $\Pr[|\bar{x}_{\ell} - x_{\ell}| > \alpha \max\{\Delta, \sigma\}] < 2 \exp(-\Omega(\alpha^2 k))$ 

Low noise ( $\sigma \leq \Delta$ ):

Same tail bound as CountSketch

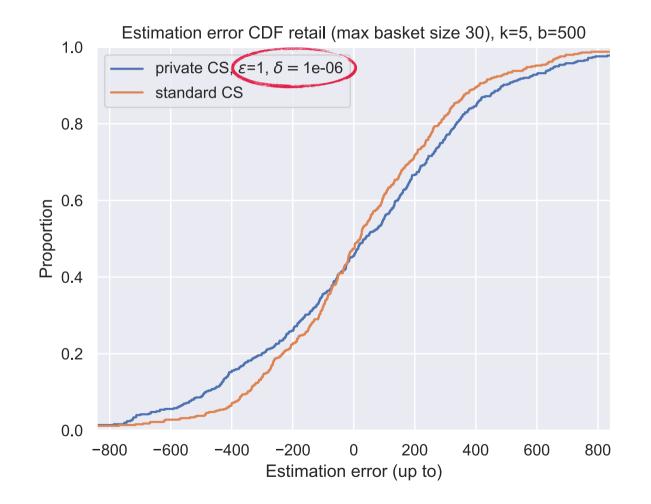
High noise ( $\sigma > \Delta$ ),  $k = \sigma^2$ : Tail like  $\mathcal{N}(0,1)$  noise +  $\exp(-\Omega(k))$ 

<u>Message of our work</u>: Estimation error of Private CountSketch is either the CountSketch error or the error needed for DP, whichever is larger

#### **Proof ingredients** (about 1 page)

- <u>Two cases</u>:
  - Adding noise with  $\sigma \leq \Delta$  maintains the probability of a good simple estimator up to a constant factor
  - Adding noise with  $\sigma > \Delta$ , the probability of a good simple estimator can be bounded up to a constant factor in terms of  $\sigma$
- Lemma from Minton & Price, using symmetry of estimators, finishes the argument

#### Experiments – market basket data



#### **Related work in NeurIPS 2022**

#### Differentially Private Linear Sketches: Efficient Implementations and Applications

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