

# Almost Optimal Algorithms for Linear Stochastic Bandits with Heavy-Tailed Payoffs

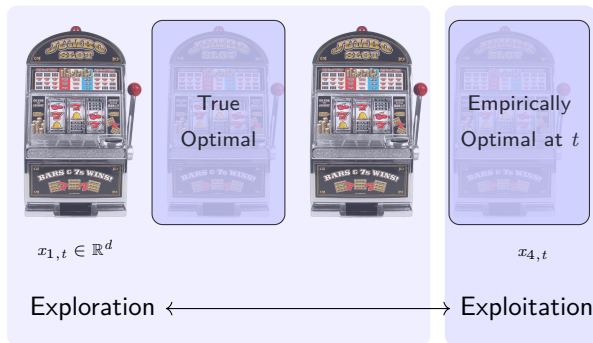
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# Linear Stochastic Bandits (LSB)

Previous setting



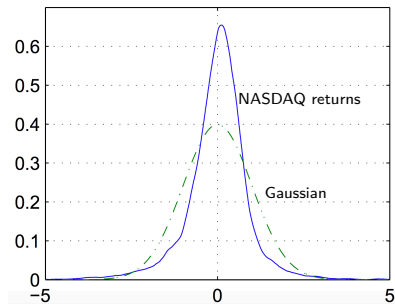
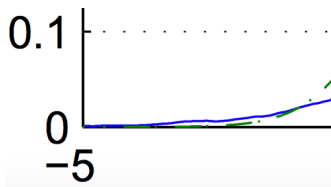
Learning setting

- ▶ 1. Given a set of arms represented by  $D \subseteq \mathbb{R}^d$
- ▶ 2. At time  $t$ , select an arm  $x_t \in D$ , and observe  $y_t(x_t) = \langle x_t, \theta_* \rangle + \eta_t$
- ▶ 3. The goal is to maximize  $\sum_{t=1}^T \mathbb{E}[y_t(x_t)]$
- ▶ 4.  $\eta_t$  follows a **sub-Gaussian** distribution ( $\mathbb{E}[\eta_t^2] < \infty$ )

# What Is A Heavy-Tailed Distribution?

## Practical scenarios

- ▶ High-probability extreme returns in financial markets



- ▶ Many other real cases
  1. Delays in communication networks (Liebeherr et al., 2012)
  2. Analysis of biological data (Burnecki et al., 2015)
  3. ...

# LSB with Heavy-Tailed Payoffs

## Problem definition

- ▶ Multi-armed bandits (MAB) with heavy-tailed payoffs (Bubeck et al., 2013)

$$\mathbb{E}[\eta_t^{1+\epsilon}] < +\infty, \quad (1)$$

where  $\epsilon \in (0, 1]$

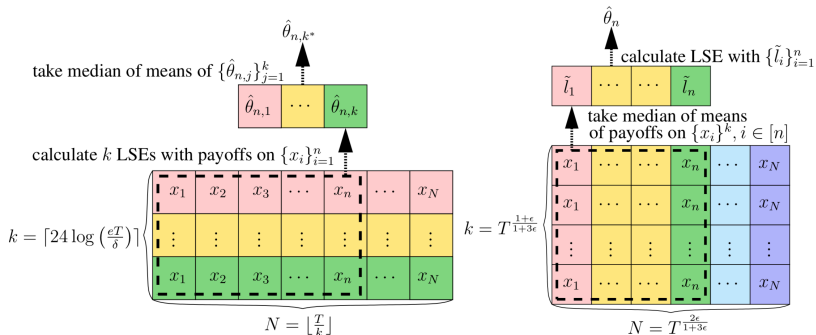
- ▶ Our setting: LSB with  $\eta_t$  satisfying Eq. (1)
  - ▶ Weaker assumption than sub-Gaussian
  - ▶ Medina and Yang (2016) studied LSB with heavy-tailed payoffs

	sub-Gaussian	heavy-tailed ( $\epsilon = 1$ )
MAB	$O(T^{\frac{1}{2}})$	$O(T^{\frac{1}{2}})$ by Bubeck et al. (2013)
LSB	$\tilde{O}(T^{\frac{1}{2}})$	$\tilde{O}(T^{\frac{3}{4}})$ by Medina and Yang (2016)

- ▶ Can we achieve  $\tilde{O}(T^{\frac{1}{2}})$ ?

# Algorithm: Median of means under OFU (MENU)

Framework comparison with MoM by Medina and Yang (2016)



(a) Framework of MENU

(b) Framework of MoM

# Regret Bounds

► Upper bounds

algorithm	MoM	MENU	CRT	TOFU
regret	$\tilde{O}(T^{\frac{1+2\epsilon}{1+3\epsilon}})$	$\tilde{O}(T^{\frac{1}{1+\epsilon}})$	$\tilde{O}(T^{\frac{1}{2} + \frac{1}{2(1+\epsilon)}})$	$\tilde{O}(T^{\frac{1}{1+\epsilon}})$

► Lower bound:  $\Omega(T^{\frac{1}{1+\epsilon}})$

When  $\epsilon = 1$ , our algorithms achieve  $\tilde{O}(T^{\frac{1}{2}})$

## See You at the Poster Session

Time: Dec. 5th, 10:45 AM – 12:45 PM

Location: Room 210 & 230 AB #158