A Unified Analysis of Stochastic Gradient Descent with Arbitrary Data Permutations and Beyond

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The Problem Formulation

(1) The finite-sum minimization problem

$$\min_{\mathbf{x} \in \mathbb{R}^d} \left[f(\mathbf{x}) \coloneqq \frac{1}{N} \sum_{n=0}^{N-1} f_n(\mathbf{x}) \right]$$

Here, N denotes the number of local objective functions, f_n denotes the local objective.

(2) Permutation-based SGD vs. classic SGD

$$\mathbf{x}^{n+1} = \mathbf{x}^n - \gamma \nabla f_{\pi(n)}(\mathbf{x}^n)$$

Here, γ denotes the step size, $\pi(n)$ is the index of the local objective at iteration n.

- Classic SGD: $\pi(n)$ is chosen uniformly with replacement from $\{0,1,...,N-1\}$.
- Permutation-based SGD: $\pi(n)$ is the (n+1)-th element of a permutation π of $\{0,1,...,N-1\}$.

The Existing Permutation-based SGD algorithms

Based on the relations among permutations, we classify the existing permutation-based SGD algorithms into the three categories:

- 1. Arbitrary Permutations (AP): Permutations are generated without any specific structure, allowing for completely arbitrary permutations for all epochs.
- 2. Independent Permutations (IP): Permutations are independent across epochs.
 - Random Reshuffling (RR): The permutation in each epoch is generated randomly.
 - FlipFlop [Rajput et al., 2022].
 - <u>Greedy Ordering</u> [Lu et al., 2022b; Mohtashami et al., 2022]: The permutation in each epoch is generated by a greedy algorithm
- 3. Dependent Permutations (DP): Permutations are dependent across epochs, with the permutation in one epoch affected by the permutations in previous epochs (explicitly).
 - <u>One Permutation (OP)</u>: The initial (first-epoch) permutation is used repeatedly for all the subsequent epochs. When the initial permutation is arbitrary, it is called <u>Incremental Gradient (IG)</u>; when the initial permutation is random, it is called <u>Shuffle Once (SO)</u>.
 - GraBs: It includes GraB [Lu et al., 2022a] and PairGraB [Cooper et al., 2023].

The Order Error

- (1) The derivation.
- For a small finite step size γ , the cumulative updates in any epoch q are

$$\mathbf{x}_{q+1} - \mathbf{x}_{q} \approx \underbrace{-\gamma N \nabla f(\mathbf{x}_{q}) + \gamma^{2} \sum_{n=0}^{N-1} \sum_{i < n} \nabla^{2} f_{\pi(n)}(\mathbf{x}_{q}) \nabla f(\mathbf{x}_{q})}_{\text{optimization vector}} + \underbrace{\gamma^{2} \sum_{n=0}^{N-1} \sum_{i < n} \nabla^{2} f_{\pi(n)}(\mathbf{x}_{q}) \left(\nabla f_{\pi(i)}(\mathbf{x}_{q}) - \nabla f(\mathbf{x}_{q}) \right)}_{\text{error vector}},$$

The goal is to suppress the error vector:

$$\|\text{Error vector}\| \leq \gamma^2 \sum_{n=0}^{N-1} \|\nabla^2 f_{\pi(n)}(\mathbf{x}_q)\| \left\| \sum_{i < n} \left(\nabla f_{\pi(i)}(\mathbf{x}_q) - \nabla f(\mathbf{x}_q) \right) \right\| \leq \gamma^2 L N \bar{\phi}_q,$$

(2) The definition.

Definition 1 (Order Error, Lu et al. [2022b,a]). The order error $\bar{\phi}_q$ in any epoch q is defined as

$$\bar{\phi}_q \coloneqq \max_{n \in [N]} \left\{ \phi_q^n \coloneqq \left\| \sum_{i=0}^{n-1} \left(\nabla f_{\pi_q(i)}(\mathbf{x}_q) - \nabla f(\mathbf{x}_q) \right) \right\|_p \right\}.$$

(3) As a measure of the quality of the permutation of examples.

The Order Error

Assumption 1 (Lu et al. [2022b,a]). There exist nonnegative constants B and D such that for all \mathbf{x}_q (the outputs of Algorithm 1),

$$(\bar{\phi}_q)^2 \le B \|\nabla f(\mathbf{x}_q)\|^2 + D.$$

Existing works implicitly deal with the order error separately across epochs (as in Asm. 1).



Assumption 2. There exist nonnegative constants $\{A_i\}_{i=1}^q$, $\{B_i\}_{i=0}^q$ and D such that for all \mathbf{x}_q (the outputs of Algorithm 1),

$$(\bar{\phi}_q)^2 \le \sum_{i=1}^q A_i (\bar{\phi}_{q-i})^2 + \sum_{i=0}^q B_i \|\nabla f(\mathbf{x}_{q-i})\|^2 + D.$$

Asm. 2 explicitly characterizes the dependence between permutations across different epochs.

In particular, when $A_i=0$, $B_i=0$ for all i=1,2,...,q, Asm. 2 reduces to Asm. 1.

Result 1: Main Theorem

Theorem 1. Let the global objective function f be L-smooth and each local objective functions f_n be $L_{2,p}$ -smooth and L_p -smooth $(p \ge 2)$. Let $\nu \ge 0$ be a numerical constant. Suppose that there exist \tilde{B} and \tilde{D} such that for $0 \le q \le \nu - 1$,

$$(\bar{\phi}_q)^2 \le \tilde{B} \|\nabla f(\mathbf{x}_q)\|^2 + \tilde{D},$$

and there exist $\{A_i\}$, $\{B_i\}$ and D such that for $q \geq \nu$,

$$(\phi_q)^2 \le \sum_{i=1}^{\nu} A_i (\bar{\phi}_{q-i})^2 + \sum_{i=0}^{\nu} B_i \|\nabla f(\mathbf{x}_{q-i})\|^2 + D.$$

The first ν epochs rely on Asm. 1, and the subsequent epochs rely on Asm. 2.

$$If \gamma \leq \min \left\{ \frac{1}{LN}, \frac{1}{32L_{2,p}N}, \frac{\sqrt{1-\sum_{i=1}^{\nu}A_{i}}}{4L_{2,p}\sqrt{\sum_{i=0}^{\nu}B_{i}}}, \frac{\sqrt{1-\sum_{i=1}^{\nu}A_{i}}}{4L_{2,p}\sqrt{\tilde{B}}}, \frac{1}{32L_{p}N} \right\}, then$$

$$Optimization term \qquad Error terms$$

$$1 \sum_{i=1}^{Q-1} \|\nabla f(z_{i})\|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} + 32^{2}L^{2} + 1 \sum_{i=1}^{Q-1} |\nabla f(z_{i})|_{2}^{2} \leq 5F_{0} +$$

It does not imposestronger constraints on γ.

where $c = \frac{10}{(1 - \sum_{i=1}^{\nu} A_i)}$ is a numerical constant.

Result 2: Case Studies

Table 7: Specific choices of A_i , B_i and D for different algorithms. The coefficients not explicitly specified equal 0. The numerical constants and polylogarithmic factors of B_i and D are omitted.

Algorithm	B_0	A_1	B_1	A_2	B_2	D	γ
AP	$N^2 \alpha^2$	0	0	0	0	$N^2 \varsigma^2$	$\gamma \lesssim \frac{1}{LN(1+\alpha)}$
RR/FlipFlop	$N^2 \alpha^2$	0	0	0	0	$N\varsigma^2$	$\gamma \lesssim \frac{1}{LN(1+\alpha/\sqrt{N})}$
GraB-proto	0	$\frac{3}{4}$ (1)	$N^2 + \alpha^2$	0	0	ς^2	$\gamma \lesssim \min\{\frac{1}{LN}, \frac{1}{L_{2,\infty}N(1+\alpha)}, \frac{1}{L_{\infty}N}\}$
GraB	0	$\frac{3}{5}$ (1)	$N^2 + \alpha^2$	$\frac{1}{50}$ (1)	N^2	ς^2	$\gamma \lesssim \min\{\frac{1}{LN}, \frac{1}{L_{2,\infty}N(1+\alpha)}, \frac{1}{L_{\infty}N}\}$
PairGraB	0	$\frac{4}{5}$ (1)	$N^2 + \alpha^2$	0	0	ς^2	$\gamma \lesssim \min\{\frac{1}{LN}, \frac{1}{L_{2,\infty}N(1+\alpha)}, \frac{1}{L_{\infty}N}\}$

A_i may take other values as long as $\sum_{i=1}^{\nu} A_i < 1$ for GraBs.

$$(\phi_q)^2 \le \sum_{i=1}^{\nu} A_i (\bar{\phi}_{q-i})^2 + \sum_{i=0}^{\nu} B_i \|\nabla f(\mathbf{x}_{q-i})\|^2 + D$$

It determines the error terms.

It does not impose stronger constraints on γ . It determines the optimization term.

Result 2: Case Studies

	Alg.	Lu et al. [2022b]	Koloskova et al. [2024] ⁽¹⁾	This work
AP ⁽²) AP	$\frac{LF_0}{Q} + \left(\frac{LF_0N\varsigma}{NQ}\right)^{\frac{2}{3}}$	$\frac{LF_0}{NQ} + \left(\frac{LF_0N\varsigma}{NQ}\right)^{\frac{2}{3}} \wedge \left(\frac{LF_0N\varsigma^2}{NQ}\right)^{\frac{1}{2}}$ (3)	$\frac{LF_0}{Q} + \left(\frac{LF_0N\varsigma}{NQ}\right)^{\frac{2}{3}}$
IP	RR	$\frac{LF_0}{Q} + \left(\frac{LF_0\sqrt{N}\varsigma}{NQ}\right)^{\frac{2}{3}}$	AP ⁽⁴⁾	$\frac{LF_0}{Q} + \left(\frac{LF_0\sqrt{N}\varsigma}{NQ}\right)^{\frac{2}{3}}$
	FlipFlop	_	_	$\frac{LF_0}{Q} + \left(\frac{LF_0\sqrt{N}\varsigma}{NQ}\right)^{\frac{2}{3}}$
DP	GraBs	_	_	$\frac{\tilde{L}F_0 + (L_{2,\infty}F_0\varsigma)^{\frac{2}{3}}}{Q} + \left(\frac{L_{2,\infty}F_0\varsigma}{NQ}\right)^{\frac{2}{3}}$ (5)

This work is the first unified framework that includes DP.

Federated Learning (Beyond Permutation-based SGD)

Algorithm 1: Permutation-based SGD

Input:
$$\pi_0$$
, \mathbf{x}_0 ; Output: $\{\mathbf{x}_q\}$

1 for $q=0,1,\ldots,Q-1$ do

2 $\mathbf{x}_q^0 \leftarrow \mathbf{x}_q$

3 for $n=0,1,\ldots,N-1$ do

4 $\mathbf{x}_q^{n+1} \leftarrow \mathbf{x}_q^n - \gamma \nabla f_{\pi_q(n)}(\mathbf{x}_q^n)$

5 $\mathbf{x}_{q+1} \leftarrow \mathbf{x}_q^N$

6 $\pi_{q+1} \leftarrow \text{Permute}(\cdots)$



We also develop a unified framework for regularized-participation FL with arbitrary permutations of clients

Algorithm 2: Regularized-participation FL

```
Input: \pi_0, \mathbf{x}_0; Output: \{\mathbf{x}_q\}
 1 for q = 0, 1, \dots, Q - 1 do
             \mathbf{w} \leftarrow \mathbf{x}_a
             for n = 0, 1, ..., N - 1 do
                    Initialize \mathbf{x}_{a,0}^n \leftarrow \mathbf{w}
                    for k = 0, 1, ..., K - 1 do
                              \mathbf{x}_{a,k+1}^n \leftarrow \mathbf{x}_{a,k}^n - \gamma \nabla f_{\pi_a(n)}(\mathbf{x}_{a,k}^n)
  6
                    \mathbf{p}_{a}^{n} \leftarrow \mathbf{x}_{a,0}^{n} - \mathbf{x}_{a,K}^{n}
                    if (n+1) \mod S = 0 then
  8
                         \mathbf{w} \leftarrow \mathbf{w} - \frac{1}{S} \sum_{s=0}^{S-1} \mathbf{p}_a^{n-s}
             \mathbf{x}_{q+1} \leftarrow \mathbf{x}_q - \eta(\mathbf{x}_q - \mathbf{w})
10
             \pi_{q+1} \leftarrow \texttt{Permute}(\cdots)
11
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Thank You!

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