

# Sublinear Time Low-Rank Approximation of Distance Matrices

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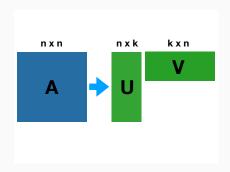
Carnegie Mellon University

- Data and Matrix compression
- De-noising and Dimensionality Reduction
- Applications to Clustering, Topic Modelling, Recommendation Systems and Distribution Learning.

#### Low-Rank Approximation

Given a  $n \times n$  matrix **A** and an integer k, compute

$$\mathbf{A}_{k} = \min_{\mathrm{rank}(\mathbf{X}) \leq k} \|\mathbf{A} - \mathbf{X}\|_{F}$$



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- 2. Let **A** be the resulting  $n \times n$  pair-wise Distance Matrix, i.e.

$$\mathbf{A} = \begin{bmatrix} \|p_1 - p_1\| & \cdot & \cdot & \|p_1 - p_n\| \\ \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot \\ \|p_n - p_1\| & \cdot & \|p_n - p_n\| \end{bmatrix}$$

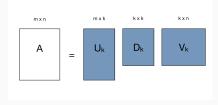
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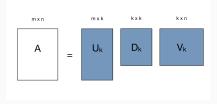
3. A is a dense matrix and has  $O(n^2)$  non-zero entries

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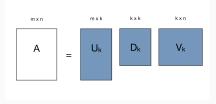


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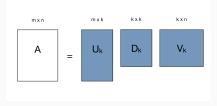
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- 3. Optimal!
- 4. Running time is  $O(n^3)$
- 5. Extremely slow for a large dataset

#### Clarkson-Woodruff showed how to output a rank k matrix **B** such that

$$\|\mathbf{A} - \mathbf{B}\|_F^2 \le (1 + \epsilon) \min_{\operatorname{rank}(\mathbf{X}) \le k} \|\mathbf{A} - \mathbf{X}\|_F^2$$

1. Running time is  $O\left(n^2 + n \operatorname{poly}\left(\frac{k}{\epsilon}\right)\right)$ 

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- 1. Running time is  $O\left(n^2 + n \operatorname{poly}\left(\frac{k}{\epsilon}\right)\right)$
- 2. Might still be too slow

Can we leverage the structure of a Distance Matrix to get faster algorithms? **Theorem :** Compute  $\mathbf{U} \in \mathbb{R}^{n \times k}$ ,  $\mathbf{V} \in \mathbb{R}^{k \times n}$  such that

$$\|\mathbf{A} - \mathbf{U}\mathbf{V}\|_F^2 \le \min_{\operatorname{rank}(\mathbf{X}) \le k} \|\mathbf{A} - \mathbf{X}\|_F^2 + \epsilon \|\mathbf{A}\|_F^2$$

in time  $O\left(n^{1.001} \operatorname{poly}\left(\frac{k}{\epsilon}\right)\right)$ 

1. Does not read most of the input!

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in time  $O\left(n^{1.001} \operatorname{poly}\left(\frac{k}{\epsilon}\right)\right)$ 

- 1. Does not read most of the input!
- 2. Only accesses  $O\left(n^{1.001} \operatorname{poly}\left(\frac{k}{\epsilon}\right)\right)$  entries in A

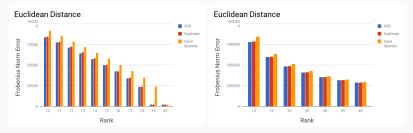
Algorithm	Running Time	
Singular Value Decomposition	O(n <sup>3</sup> )	
Input Sparsity Low-Rank Approximation	$O\left(n^2 + n \operatorname{poly}\left(\frac{k}{\epsilon}\right)\right)$	
Sublinear Low-Rank Approximation	$O\left(n^{1.001} \operatorname{poly}\left(\frac{k}{\epsilon}\right)\right)$	

Algorithm	Clustering	MNIST
Singular Value Decomposition	398.76	398.50
Input Sparsity Low-Rank Approximation	8.94	34.32
Sublinear Low-Rank Approximation	1.69	4.16

#### **Experiments: Absolute Error**



#### **MNIST** Dataset



**Figure 1:** We plot  $||\mathbf{A} - \mathbf{B}||_F$  on a synthetic dataset with 20 clusters and the MNIST dataset using  $\ell_2$  as the metric. We compare the error achieved by SVD (optimal), our Sublinear Algorithm and the Input Sparsity Algorithm.

Thank You!